

# Distributed Authorization Framework for Mobile Agents

**G. Navarro**<sup>1</sup>   J. A. Ortega-Ruiz<sup>2</sup>   J. Ametller<sup>1</sup>   S. Robles<sup>1</sup>

<sup>1</sup>SeNDA *Security for Network and Distributed Applications* Group  
Dept. of Information and Communications Engineering  
Universitat Autònoma de Barcelona, Spain.

<sup>2</sup>Institute for Space Studies of Catalonia (IEEC),  
80034 Barcelona, Spain.

MATA 2005, Montreal.

# Outline

- 1 Introduction
- 2 Naming Schema
  - Local Names
  - Groups/roles
- 3 Authorization
- 4 DAFMA Components
  - AM and RM
  - RC and CRM
  - Simple Architecture example
  - Establishing MA's role membership
- 5 Conclusions

# Introduction

## DAFMA

### Distributed **A**uthorization **F**ramework for **M**obile **A**gents

- Provides an authorization framework for multiagent systems supporting agent mobility.
- Based on **trust management** systems.
- Provides Role-Based Access Control (RBAC) services.

# Introduction (II)

- Agent mobility introduces limitations. Mobile agents:
  - May be able to directly perform cryptographic operations?
  - May have cryptographic keys?

# Introduction (II)

- Agent mobility introduces limitations. Mobile agents:
  - May be able to directly perform cryptographic operations?
  - May have cryptographic keys?
- MAs in DAFMA do not need to perform cryptographic operations.

# Naming schema

- Strongly influenced by SPKI/SDSI's **local names**.

# Naming schema

- Strongly influenced by SPKI/SDSI's **local names**.
- Each principal, is identified by:
  - Public key: non-MA principals. They have a cryptographic key pair.
  - Hash of an object: MA principals.

# Local names

- Each entity as a **name container** for defining local names.

(`<principal>`, `<local-name>`)

# Local names

- Each entity as a **name container** for defining local names.

(`<principal>`, `<local-name>`)

- Example:  $PK_0$  defines a local name *partner* for  $PK_1$ :

$PK_0: (PK_1, partner)$

# Local names

- Each entity as a **name container** for defining local names.

(`<principal>`, `<local-name>`)

- Example:  $PK_0$  defines a local name *partner* for  $PK_1$ :

$PK_0$ : ( $PK_1$ , partner)

- Third parties can use a *fully qualified name* to refer to  $PK_0$ 's *partner*.

$PK_0$  partner

# Local names

- Each entity as a **name container** for defining local names.

(`<principal>`, `<local-name>`)

- Example:  $PK_0$  defines a local name *partner* for  $PK_1$ :

$PK_0: (PK_1, \text{partner})$

- Third parties can use a *fully qualified name* to refer to  $PK_0$ 's *partner*.

$PK_0$  partner

- Name assertion: local names can be made public through a signed statement  $\rightarrow$  **name certificate**:

$\{(PK_1, \text{partner})\}_{PK_0^{-1}}$

# Group/role Management

- Groups/roles created as local names:

$$\{(PK_1, \text{friends})\}_{PK_{adm}^{-1}}$$
$$\{(PK_2, \text{friends})\}_{PK_{adm}^{-1}}$$
$$\{(PK_0 \text{ partner}, \text{friends})\}_{PK_{adm}^{-1}}$$

# Group/role Management

- Groups/roles created as local names:

$$\{(PK_1, \text{friends})\}_{PK_{adm}^{-1}}$$
$$\{(PK_2, \text{friends})\}_{PK_{adm}^{-1}}$$
$$\{(PK_0 \text{ partner}, \text{friends})\}_{PK_{adm}^{-1}}$$

- Support for group/role **hierarchies** by means of inclusion.  
→ support for RBAC-like systems.

$$\{(PK_{adm} \text{ family}, \text{friends})\}_{PK_{adm}^{-1}}$$

# Authorization certificates

- Authorization expressed through **authorization certificates** (SPKI/SDSI).

# Authorization certificates

- Authorization expressed through **authorization certificates** (SPKI/SDSI).
- Includes a boolean delegation statement: determines if the authorization may be further delegated or not.

# Authorization certificates

- Authorization expressed through **authorization certificates** (SPKI/SDSI).
- Includes a boolean delegation statement: determines if the authorization may be further delegated or not.
- Delegation adds a lot of flexibility to the system:
  - But introduces complexity: we need reduction and resolution algorithms.
  - Currently used a simplified version of SPKI/SDSI's certificate chain discovery algorithms.

# DAFMA Architecture

- Divided in 4 components:
  - AM** Authorization Manager.
  - RM** Role Manager.
  - RC** Resource Controller.
  - CRM** Certificate Repository Manager.

# DAFMA Architecture

- Divided in 4 components:
  - AM** Authorization Manager.
  - RM** Role Manager.
  - RC** Resource Controller.
  - CRM** Certificate Repository Manager.
- Implemented as static agents.

# DAFMA Architecture

- Divided in 4 components:
  - AM** Authorization Manager.
  - RM** Role Manager.
  - RC** Resource Controller.
  - CRM** Certificate Repository Manager.
- Implemented as static agents.
- Interaction between components: FIPA Agent Communication Language and Ontologies.

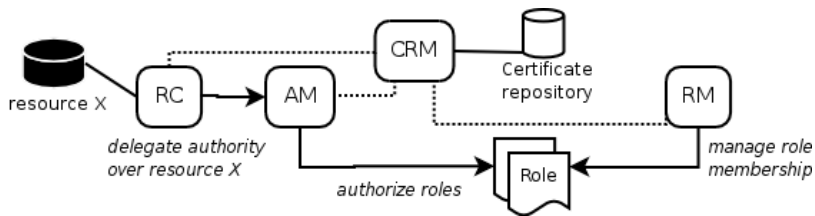
# AM and RM

- Authorization Manager:
  - **Authorization** → **roles**.
  - Follows local authorization policy.
  
- Role Manager:
  - **Role membership**.
  - Follows local role policy.

# RC and CRM

- Resource Controller:
  - Acts like a Policy Enforcement Point.
  - Simple authorization policy: determines AMs allowed to manage authorizations for the resource.
- Certificate Repository Manager:
  - Certificate repository and related services.

# Simple DAFMA Architecture



# Establishing MA's role membership

- Two supported role membership establishment:
  - **User-managed** role:
    - RM gives full control to a user to manage its role membership.
  - **RM-managed** role:
    - The users needs to query the RM each role membership.

# Conclusions

- DAFMA provides an authorization system for distributed environments taking into account agent mobility.

# Conclusions

- DAFMA provides an authorization system for distributed environments taking into account agent mobility.
- Mobile agents do not need to perform cryptographic operations nor need to carry sensitive information (private keys).

# Conclusions

- DAFMA provides an authorization system for distributed environments taking into account agent mobility.
- Mobile agents do not need to perform cryptographic operations nor need to carry sensitive information (private keys).
- Based on trust management and RBAC.

# Conclusions

- DAFMA provides an authorization system for distributed environments taking into account agent mobility.
- Mobile agents do not need to perform cryptographic operations nor need to carry sensitive information (private keys).
- Based on trust management and RBAC.
- Current prototype implemented on top of JADE, supporting FIPA specifications.

# Thanks

Thanks. Questions?